gputils 0.13.0

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Chapter 1

Introduction

gputils is a collection of tools for Microchip (TM) PIC microcontrollers. It includes gpal, gpasm, gplink, and gplib. Each tool is intended to be an open source replacement for a corresponding Microchip (TM) tool. This manual covers the basics of running the tools. For more details on a microcontroller, consult the manual for the specific PICmicro product that you are using.

This document is part of gputils.

1.1 Tool Flows

gputils can be used in three different ways: absolute asm mode, relocatable asm mode, and HLL mode.

1.1.1 Absolute Asm Mode

In absolute asm mode, an assembly language source file is directly converted into a hex file by gpasm. This method is absolute because the final addresses are hard coded into the source file.

1.1.2 Relocatable Asm Mode

In relocatable asm mode, the microcontroller assembly source code is divided into separate modules. Each module is assembled into an object using gpasm. That object can be placed “anywhere” in microcontroller’s memory. Then gplink is used to resolve symbols references, assign final address, and to patch the machine code with the final addresses. The output from gplink is an absolute executable object.

1.1.3 HLL Mode

In HLL (High Level Language) mode, the source code is written in a Ada like language. gpal then converts that file into a relocatable assembly file. It then automatically invokes gpasm and gplink to generate an absolute executable object.

1.1.4 Which Tool Flow is best?

Absolute mode is simple to understand and to use. It only requires one tool, gpasm. Most of the examples on Microchip's website use absolute mode. So why use relocatable mode?

- Code can be written without regard to addresses. This makes it easier to write and reuse.
- The objects can be archived to create a library, which also simplifies reuse.
- Recompiling a project can be faster, because you only compile the portions that have changed.
- Files can have local name spaces. The user chooses what symbols are global.

Most develpment tools use relocatable objects for these reasons. The few that don't are generally microcontroller tools. Their applications are so small that absolute mode isn't impractical. For PICs, relocatable mode has one big disadvantage. The bank and page control is a challenge. To overcome that, HLL mode can be used. It helps to hide these details from the user.

1.2 Supported processors

gputils currently supports the following processors:

eeprom8	gen	p10f200	p10f202	p10f204	p10f206
p12c508	p12c508a	p12c509	p12c509a	p12c671	p12c672
p12ce518	p12ce519	p12ce673	p12ce674	p12cr509a	p12f508
p12f509	p12f629	p12f635	p12f675	p12f683	p14000
p16c5x	p16cxx	p16c432	p16c433	p16c505	p16c52
p16c54	p16c54a	p16c54b	p16c54c	p16c55	p16c55a
p16c554	p16c557	p16c558	p16c56	p16c56a	p16c57
p16c57c	p16c58a	p16c58b	p16c61	p16c62	p16c62a
p16c62b	p16c620	p16c620a	p16c621	p16c621a	p16c622
p16c622a	p16c63	p16c63a	p16c64	p16c64a	p16c642
p16c65	p16c65a	p16c65b	p16c66	p16c662	p16c67
p16c71	p16c710	p16c711	p16c712	p16c715	p16c716
p16c717	p16c72	p16c72a	p16c73	p16c73a	p16c73b
p16c74	p16c745	p16c747	p16c74a	p16c74b	p16c76
p16c765	p16c77	p16c770	p16c771	p16c773	p16c774
p16c781	p16c782	p16c84	p16c923	p16c924	p16c925
p16c926	p16ce623	p16ce624	p16ce625	p16cr54	p16cr54a
p16cr54b	p16cr54c	p16cr56a	p16cr57a	p16cr57b	p16cr57c
p16cr58a	p16cr58b	p16cr62	p16cr620a	p16cr63	p16cr64
p16cr65	p16cr72	p16cr83	p16cr84	p16f505	p16f54
p16f57	p16f59	p16f627	p16f627a	p16f628	p16f628a

p16f630	p16f636	p16f639	p16f648a	p16f676	p16f684
p16f685	p16f687	p16f688	p16f689	p16f690	p16f716
p16f72	p16f73	p16f737	p16f74	p16f76	p16f767
p16f77	p16f777	p16f785	p16f818	p16f819	p16f83
p16f84	p16f84a	p16f87	p16f870	p16f871	p16f872
p16f873	p16f873a	p16f874	p16f874a	p16f876	p16f876a
p16f877	p16f877a	p16f88	p16f913	p16f914	p16f916
p16f917	p16hv540	p17cxx	p17c42	p17c42a	p17c43
p17c44	p17c752	p17c756	p17c756a	p17c762	p17c766
p17cr42	p17cr43	p18cxx	p18c242	p18c252	p18c442
p18c452	p18c601	p18c658	p18c801	p18c858	p18f1220
p18f1320	p18f2220	p18f2320	p18f2331	p18f2410	p18f242
p18f2420	p18f2431	p18f2439	p18f2455	p18f248	p18f2480
p18f2510	p18f2515	p18f252	p18f2520	p18f2525	p18f2539
p18f2550	p18f258	p18f2580	p18f2585	p18f2610	p18f2620
p18f2680	p18f2681	p18f4220	p18f4320	p18f4331	p18f4410
p18f442	p18f4420	p18f4431	p18f4439	p18f4455	p18f448
p18f4480	p18f4510	p18f4515	p18f452	p18f4520	p18f4525
p18f4539	p18f4550	p18f458	p18f4580	p18f4585	p18f4610
p18f4620	p18f4680	p18f4681	p18f6310	p18f6390	p18f6410
p18f6490	p18f64j15	p18f6520	p18f6525	p18f6585	p18f65j10
p18f65j15	p18f6620	p18f6621	p18f6627	p18f6680	p18f66j10
p18f66j15	p18f6720	p18f6722	p18f67j10	p18f8310	p18f8390
p18f8410	p18f8490	p18f84j15	p18f8520	p18f8525	p18f8585
p18f85j10	p18f85j15	p18f8620	p18f8621	p18f8627	p18f8680
p18f86j10	p18f86j15	p18f8720	p18f8722	p18f87j10	rf509af
rf509ag	rf675f	rf675h	rf675k	sx18	sx20
sx28					

Chapter 2

gpal

2.1 Introduction

gpal is a compiler for Microchip (TM) PIC microcontrollers. Unlike most of the other tools in gputils, there is no corresponding Microchip tool that it replaces. It is a new tool and language specifically designed to to simplify software development for PICs. The language to very similar to the Pascal family of languages, specifically Ada.

gpal was inspired by Jal <http://jal.sourceforge.net>. That language was created by Wouter Van Ooijen.

2.2 Running gpal

The general syntax for running gpal is

```
gpal [options] input-files
```

Where options can be one of:

Option	Meaning
a	Compile or assemble, then archive
c	Compile or assemble, but don't link
d	Output debug messages
h	Show the usage message
H	Scan the specified processor header file
I <directory>	Specify an include directory
k"<options>"	Extra link or lib options
l	List supported processors
M	Generate a Make compatible dependency list
o <file>	Alternate name of hex output file
O<level>	Optimization level
p<processor>	Select target processor
q	Quiet
S	Compile only, don't assemble or link
t	Do not delete intermediate files
v	Print gpal version information and exit

2.2.1 Operations

gpal only converts .pal source files into .asm files. However, as a convience it can automatically invoke gpasm to convert the .asm file into an object file with a .o extension. It can also invoke gplink to produce a PIC executable or gplib to produce an archive of objects. The operations are selected using the options -S, -a, and -c.

gpal will automatically remove any temporary file generated by its operation or by an tool it invokes. That behavior can be controlled using the -t option.

Currently there is no difference between invoking gpal with a complete list of input files, verses invoking multiple times, once for each file. In the future, that will probably change.

2.2.2 Input files

gpal will attempt to compile any input file regardless of its name or extension. Typically two file extensions, .pal and pub, are used.

2.3 Basics

2.3.1 Free-format

So this statement:

```
if a>b then
  timer = 0;
end if;
```

is equivalent to:

```
if a>b then timer = 0; end if;
```

although not recommended.

2.3.2 Statement terminator

The semicolon is used to terminate all statements and subprograms.

2.3.3 Comments

Comments are proceeded by a double minus (–) and continue until the end of the current line.

2.4 Types

2.4.1 Builtin Types

The following table defines the built in types:

Name	Size in bytes	Minimum Value	Maximum Value
uint8	1	0	255
int8	1	-128	127
uint16	2	0	65,535
int16	2	-32,768	32,767
uint24	3	0	16,777,215
int24	3	-8,388,608	8,388,607
uint32	4	0	4,294,967,295
int32	4	-2,147,483,648	2,147,483,647

None of the ranges are checked at run time. The user must ensure that any assignment expression won't overflow or underflow the type.

2.4.2 Access types

```
type <name> is access <type>;
```

Will create an access type to <type>. This will allow indirect access to any variable of type <type>.

2.4.3 Arrays

```
type <name> is array <expression> to <expression> of <type>;
```

The following example will create an array type of 10 unsigned bytes:

```
type buffer_type is array 1 to 10 of uint8;
variable buffer : buffer_type;
```

2.4.4 Enumerated

```
type <name> is ( <name> [, <name>]*);
```

The following code will create and enumerated type:

```
type main_state is (INIT, DELAY, OUTPUT);
```

This will create a new type that can take on one of three values. Each symbol in the list is assigned a value starting at 0. Each symbol value pair is added to the global symbol table. All enumerated types use the uint8 size. So there for the maximum list size is 256 members.

2.4.5 Records

```
type <name> is record
  [<name> : <type>;]*
end record;
```

A record is an aggregate of declarations. It allows complex variable structures to be defined with a single statement. It also allows that definition to be easily accessed.

```
type mode_type is (IDLE, ACTIVE, ERROR);
type control_type is record
  time : uint8;
  mode : mode_type;
end record;
```

A variable could then be defined like this:

```
control : control_type = (0, IDLE);
```

This variable “control” would be accessed in the body of a subprogram as follows:

```
control.time = 10;
control.mode = ACTIVE;
```

2.4.6 Type Alias

Types can be given new names to suit the user’s preference.

```
type <name> is <type>;
```

This example will create an alias of int16 with the name short.

```
type short is int16;
```

2.5 Expressions

2.5.1 Symbols

Symbols must match the following rule:

```
[_.a-z][_.0-9a-z]*
```

All symbols are case insensitive. So the following two statements are equivalent.

```
Timer  = 0;
tImEr  = 0;
```

The only exception is symbols that used to generate filenames.

```
with  time;
```

This statement will open the file “time.pub”. If the host operating system uses a case sensitive file system, the case of the with is important. To maintain portability across different operating systems, it is best to keep the with statements and filenames lower case.

2.5.2 Attributes

Attributes provide a means to access the properties of symbols. They always use the tick symbol (’). They are in the form:

```
<symbol>'<attribute>
```

Access

Indirectly access memory using the contents of the symbol as the address. This is only valid on symbols that are declared as access types. Using it on any other symbol will cause an error to be generated.

Address

Generates relocation symbols for the address of the variable. Note that the address is generally not known at compile time. It will be patched by the linker with the correct value after relocation.

First

Returns the index of the first element in an array. It is only valid on array symbols.

Last

Returns the index of the last element in an array. It is only valid on array symbols.

Range

Equivalent to:

```
<symbol>'first TO <symbol>'last
```

Only available in the range statement of a FOR loop.

Size

Returns the size of a variable in bits.

2.5.3 Numbers

gpal uses decimal as its default radix. The following table summarizes other supported numeric formats.

base	general syntax	21 decimal written as
decimal	[0-9]*	21
hex	0x[0-F]*	0x15

2.5.4 Operators

gpal supports a full set of operators, based on the C operator set. The operators in the following table are arranged in groups of equal precedence, but the groups are arranged in order of increasing precedence. When gpal encounters operators of equal precedence, it always evaluates from left to right.

Operator	Description
=	assignment
||	logical or
&&	logical and
&	bitwise and
|	bitwise or
^	bitwise exclusive-or
<	less than
>	greater than
==	equals
!=	not equals
>=	greater than or equal
<=	less than or equal
<<	left shift
>>	right shift
+	addition
-	subtraction
*	multiplication
/	division
%	modulo
-	negation
!	logical not
~	bitwise no

2.5.5 Assignment

<name>['[' <expression> ']']? = <expression>;

Assignment statements can appear in any statement block. <name> must be a variable. If the bracket enclosed expression is added it must be an array.

2.5.6 Test

<expression> [<comparison operator> <expression>]*;

Test statements can only appear in the expressions of if statements and while loops. They must evaluate to a boolean.

2.5.7 Label

Labels must match the following rule:

[_.a-z][_.0-9a-z]*:

All labels are case insensitive. So the following two statements are equivalent.

```
MY_LaBeL:
my_lAbEl:
```

Labels can only appear in the body of a subprogram. Labels are only valid within the subprogram that they appear.

2.6 Statements

2.6.1 Assembly

```
asm
  <asm statements>
end asm;
```

Unmodified <asm statements> are copied to the assembly file output. The syntax of <asm statements> must be compatible with gpasm.

2.6.2 Case

```
case <name> is
[when <constant>  [ | <constant>  ]* =>
  <statements>]*
[when others =>
  <statements>]?
end case;
```

If <name> equals any of the <constants> the <statements> are executed. If none of the constants match and an others is present, the others statements are executed. Here is an example:

```
case input is
  when MAXIMUM | MINIMUM =>
    output = 8;
  when 5 | 3 | 2 =>
    output = 4;
  when 1 =>
    output = 2;
  when others =>
    output = 0;
end case;
```

2.6.3 For

```
for <name> in <start_expression>  to <end_expression>
  loop
    <statements>
  end loop;
```

<name> is set to <start_expression>. It is then incremented each time the block of statements are executed. It continues until <name> reaches <end_expression>. Here is an example:

```
for i in 0 to 10
  loop
    buffer[i]  = 0;
  end loop;
```

2.6.4 Goto

```
goto <label>;
```

Branch progam flow to the address specified by <label>. <label> must occur with in the same subprogram as the goto. <label> can appear before or after the goto statement.

The use of goto is discouraged in most languages, but most have them. They can aid in the creation of complex state machines. For gpal, it is the mechanism for exiting infinite loops or exiting finite loops early.

```
loop
  j = j + 1;
  if j = 100 then
    goto end_of_loop;
  end if;
end loop;
end_of_loop:
```

2.6.5 If

```
if <expression>  then
  <statements>
[elsif <expression>  then
  <statements>]*
[else
  <statements>]?
end if;
```

The statements in each block are executed if the expression is true. Here is an example:

```
if i < 10 then
  j = 5;
elsif i > 12 then
  j = 10;
elsif i > 14 then
  j = 14;
else
  j = 0;
end if;
```

2.6.6 Loop

```
loop
  <statements>
end loop;
```

The statements in the block are executed in an infinite loop. Here is an example:

```
loop
  j = j + 1;
  if j = 100 then
    return 0;
  end if;
end loop;
```

2.6.7 Null

```
null;
```

Execute a NOP. This is a little different from NULL statements in most languages. It is typically is used in statement block that was intentionally left blank and no code is generated.

2.6.8 Pragma

```
pragma <anything>;
```

Pragmas provide data to compiler which is outside of its legal syntax. The table below summarizes the pragmas available:

Name	Format	Description
Code Address	code_address = <constant>	Make the code section absolute at address <constant>.
Code Section	code_section = “<name>“	Set the code section name to <name>.
Processor	processor = “<name>“	Set the processor name to <name>.
Data Address	data_address = <constant>	Make the data section absolute at address <constant>.
Data Section	data_section= “<name>“	Set the data section name to <name>.

2.6.9 Return

```
return <expression>;
```

Evaluate the <expression>, place it in the return register, and return from the function.

2.6.10 While

```
while <expression>
  loop
    <statements>
  end loop;
```

The statements in the block are executed while <expression> is true. Here is an example:

```
while j < 10
  loop
    j = j + 1;
  end loop;
```

2.6.11 With

```
with <name>;
```

The with statement tells the compiler to add the data from the public <name> to its symbol tables. This will allow access to that module's subprograms and data.

2.7 Declarations

2.7.1 Variables

Symbols whose values change during runtime are referred to as variables. Because variables change value during run time they are stored in data memory. An expression specifies its initial value and the address of the variable. A variable is declared as follows:

```
<name> : <type> [= <expression>]? [at <expression>]?;
```

Here is an example:

```
gain : short = 10 at 0x30;
```

As shown in the example above, the address of the variable can be specified. This feature is not available for variables declared within a subprogram. This is available for global data. However, this should generally be avoided. Manually assigning addresses can interfere with optimial relocation of the data memory sections. This could result in more bank switching and a more fragmented memory map.

2.7.2 Constants

Compile time symbols whose values do not change are referred to as constants. A constant is declared as follows:

```
<name> : constant = <expression>;
```

Here is an example:

```
filter_offset : constant = 0x1434;
```

2.7.3 Alias

An alias can be created for complex expressions that that used often.

```
alias <alias name> <expression>;
```

This example will create an alias of p16f877.porta with the name vdata.

```
alias vdata p16f877.porta;
```

2.8 Subprograms

2.8.1 Procedure

```
procedure <name> ( [<arg name> : [in|out|inout] <type>]* )? is
  <declarations>
begin
  <statements>
end procedure;
```

This creates a block of executable code that starts at <name>. The procedure can be called from other subprograms within any statement block, but they can not be called from within an expression.

Permanent storage is allocated for each procedure argument. Data is passed to and from the procedure through that storage. The calling subprogram puts data into the arguments and reads from the arguments based the direction specified in the procedure definition. The direction is ignored by the procedure. All the arguments can be read from and written to.

Local constants and variables are declared in <declarations>. Any variables declared in this region may be permanent or overlayed with data from other subprograms.

2.8.2 Function

```
function <name> ( [<arg name> : [in|out|inout] <type>]* )? return type is
  <declarations>
begin
  <statements>
end function;
```

This creates a block of executable code that starts at <name>. The function can only be called from within expressions.

Like procedures in many respects, except a value is returned. This value is used in the expression.

2.9 Files

Input files are text only. They can have any name or extension. The files are composed of modules and publics. Any number of modules or publics can be in a single file.

2.9.1 Module

```
module <name> is
  <subprogram definitions|variable definitions|constants|types>
end module;
```

The module defines a related group of subprograms and data that will be placed in the same page or bank. Groups of modules are compiled and linked to gather to make the executable.

Each module will be written to one assembly file and subsequently result in one object file to be generated. Typically, one module is placed in each .pal file. The filename will be the same as <name> with the .pal extension added.

2.9.2 Public

```
public <name> is
  <subprograms declarations|variable declarations|constants|types>
end public;
```

The public declares which portions of its module will be visible to other modules. It also provides information about the interface to the module's subprograms and data. Typically, one public is placed in each .pub file. The filename will be the same as <name> with the .pub extension added. The <name> must also match the <name> of its module if one exists. When the module is compiled, it will scan its public file to verify that the declarations in the file match its subprograms and data.

2.10 Vectors

2.10.1 Reset

A reset vector and the associated startup code is generated on any procedure with the name "main". It can reside in any module, but there can be only one procedure with the name "main".

2.10.2 Interrupt

An interrupt vector and the associated context saving code is generated on any procedure with the name "isr". It can reside in any module, but there can be only one procedure with the name "isr".

2.11 Code Generation

2.11.1 Phases

Parse

The input files are parsed and stored in memory in a tree format. Constructs are replaced with common structures. For example, for and while loops are converted into conditional loops with initialization statements and increment statements.

Analyze

The syntax and semantics of the tree are checked. Most of the errors are generated during this phase.

Optimize Tree

The tree is modified to generate better code. This is a high level optimization.

Code Generation

The icode is written to an asm file, so it can be assembled and linked.

2.11.2 Expression Evaluation

Most compilers are a stack machine, accumulator machine, or a register machine. Not all of the machines are good for every target processor. Some options either won't work or aren't optimal. Because of the limited resources on PICs, the absence of stack manipulation instructions, and the fact that some instructions can only target the Wreg. An accumulator machine is the best choice for PICs. It is the choice that gpal uses. For example:

```
i = (x + 3) & 4;
```

will generate the following pseudo code:

```
Wreg = x;
Wreg = Wreg + 3;
Wreg = Wreg & 4;
i = Wreg;
```

If necessary, intermediate values are stored in data memory. For byte sized operations the Wreg is used. For larger sized operations, a section of data memory is used as the accumulator.

2.11.3 COFF sections

All the executable code in a module is placed in one COFF section. This guarantees that code will be on the same page, so no page switching is required. the code section name is specified using the code_section pragma. If that pragma isn't used a default name is used.

Similarly all data memory is placed in one COFF section. This too reduces the number of bank switches when accessing local data.

2.11.4 Name mangling

gpal uses a hierarchical name space for all symbols and types. The module name is specified in all inter-module accesses. So to write to memory "clock" in the local module:

```
clock = 0;
```

to write to memory "clock" in module "time":

```
time.clock    = 0;
```

To prevent collisions when compiling and linking all symbols are mangled in the asm output of gpal. It takes the form:

```
<module>.<subprogram>.<local          data>
```

So, local data "index" in procedure "pop_stack" in module "stack" is given the name "stack.pop_stack.index".

2.12 Coding Suggestions

The following suggestions will help to generate smaller faster target code.

2.12.1 Use uint8 types

PICs are unsigned 8 bit machines. To do anything beyond that requires more memory and more instruction cycles. So use uint8 for as many arguments and data as possible.

2.12.2 Keep data private

Any subprogram or data in the public will make the corresponding object public. When an object is public fewer compile time optimizations can be done. For example, if public data is defined in a module, but not used in that module, it can't be removed. Another module may access that data. So put as few subprogram and data declarations in the public file as possible.

2.12.3 Group related subprograms and data in one module

Any time data is accessed in another module, it could require a bank switch. Minimizing switches will help to reduce code size and increase speed.

2.12.4 Name COFF sections

Before relocating sections, gplink will combine all like named sections into one larger section. Sections can not cross page or bank boundaries, so inter section accesses don't require bank or page switches. To name the sections use the code_section and udata_section pragmas. This will group the code together. If the code is needed on a specific page or bank, create a logical definition in your linker script.

2.12.5 Don' t use absolute sections

gpal provides the ability to specify the address of the code or udata of the current module. The feature is provided for a limited set of cases were the address must be known. Unfortunately, it limits the choices the linker can make when relocating sections. At worst it may make the design not fit in the available memory. It also requires extra effort on the part of the user.

2.12.6 Use multiple module implementations

The public file defines the interface to a module. The name of the public and its file must match, so it can be found. There is no requirement of module name to match the file name. This makes it possible for multiple modules with the same interface to exist. You select the module to use when the project is linked. For example you could have math_fast.pal and math_small.pal that both contain a math module. One written to execute fast and the other written to be small in memory.

Chapter 3

gpasm

3.1 Running gpasm

The general syntax for running gpasm is

```
gpasm [options] asm-file
```

Where options can be one of:

Option	Meaning
a <format>	Produce hex file in one of four formats: inhx8m, inhx8s, inhx16, inhx32 (the default).
c	Output a relocatable object
d	Output debug messages
D symbol[=value]	Equivalent to "#define <symbol> <value>"
e [ON|OFF]	Expand macros in listing file
g	Use debug directives for COFF
h	Display the help message
i	Ignore case in source code. By default gpasms to treats "fooYa" and "FOOYA" as being different.
I <directory>	Specify an include directory
l	List the supported processors
L	Ignore nolist directives
m	Memory dump
n	Use DOS style newlines (CRLF) in hex file. This option is disabled on win32 systems.
o <file>	Alternate name of hex output file
p<processor>	Select target processor
q	Quiet
r <radix>	Set the radix, i.e. the number base that gpasm uses when interpreting numbers.<radix> can be one of "oct", "dec" and "hex" for bases eight, ten, and sixteen respectively. Default is "hex".
v	Print gpasm version information and exit
w [0 | 1 | 2]	Set the message level
y	Enable 18xx extended mode

Unless otherwise specified, gpasm removes the ".asm" suffix from its input file, replacing it with ".lst" and ".hex" for the list and hex output files respectively. On most modern operating systems case is significant in filenames. For this reason you should ensure that filenames are named consistently, and that the ".asm" suffix on any source file is in lower case.

gpasm always produces a ".lst" file. If it runs without errors, it also produces a ".hex" file or a ".o" file.

3.1.1 Using gpasm with "make"

On most operating systems, you can build a project using the make utility. To use gpasm with make, you might have a "makefile" like this:

```
tree.hex:  tree.asm  treedef.inc
        gpasm  tree.asm
```

This will rebuild "tree.hex" whenever either of the "tree.asm" or "treedef.inc" files change. A more comprehensive example of using gpasm with makefiles is included as example1 in the gpasm source distribution.

3.1.2 Dealing with errors

gpasm doesn't specifically create an error file. This can be a problem if you want to keep a record of errors, or if your assembly produces so many errors that they scroll off the screen. To deal with this if your shell is "sh", "bash" or "ksh", you can do something like:

```
gpasm tree.asm 2>&1 | tee tree.err
```

This redirects standard error to standard output ("2>&1"), then pipes this output into "tee", which copies it input to "tree.err", and then displays it.

3.2 Syntax

3.2.1 File structure

gpasm source files consist of a series of lines. Lines can contain a label (starting in column 1) or an operation (starting in any column after 1), both, or neither. Comments follow a ";" character, and are treated as a newline. Labels may be any series of the letters A-z, digits 0-9, and the underscore ("_"); they may not begin with a digit. Labels may be followed by a colon (":").

An operation is a single identifier (the same rules as for a label above) followed by a space, and a comma-separated list of parameters. For example, the following are all legal source lines:

```
                              ; Blank line
loop    sleep                 ; Label and operation
        incf    6,1           ; Operation with 2 parameters
        goto    loop          ; Operation with 1 parameter
```

3.2.2 Expressions

gpasm supports a full set of operators, based on the C operator set. The operators in the following table are arranged in groups of equal precedence, but the groups are arranged in order of increasing precedence. When gpasm encounters operators of equal precedence, it always evaluates from left to right.

Operator	Description
=	assignment
||	logical or
&&	logical and
&	bitwise and
|	bitwise or
^	bitwise exclusive-or
<	less than
>	greater than
==	equals
!=	not equals
>=	greater than or equal
<=	less than or equal
<<	left shift
>>	right shift
+	addition
-	subtraction
*	multiplication
/	division
%	modulo
UPPER	upper byte
HIGH	high byte
LOW	low byte
-	negation
!	logical not
~	bitwise no

Any symbol appearing in column 1 may be assigned a value using the assignment operator (=) in the previous table. Additionally, any value previously assigned may be modified using one of the operators in the table below. Each of these operators evaluates the current value of the symbol and then assigns a new value based on the operator.

Operator	Description
=	assignment
++	increment by 1
–	decrement by 1
+=	increment
-=	decrement
*=	multiply
/=	divide
%=	modulo
<<=	left shift
>>=	right shift
&=	bitwise and
|=	bitwise or
^=	bitwise exclusive-or

3.2.3 Numbers

gpasm gives you several ways of specifying numbers. You can use a syntax that uses an initial character to indicate the number's base. The following table summarizes the alternatives. Note the C-style option for specifying hexadecimal numbers.

base	general syntax	21 decimal written as
binary	B'[01]*'	B'10101'
octal	O'[0-7]*'	O'25'
decimal	D'[0-9]*'	D'21'
hex	H'[0-F]*'	H'15'
hex	0x[0-F]*	0x15

When you write a number without a specifying prefix such as "45", gpasm uses the current radix (base) to interpret the number. You can change this radix with the RADIX directive, or with the "-r" option on gpasm's command-line. The default radix is hexadecimal.

If you do not start hexadecimal numbers with a digit, gpasm will attempt to interpret what you've written as an identifier. For example, instead of writing C2, write either 0C2, 0xC2 or H'C2'.

Case is not significant when interpreting numbers: 0ca, 0CA, h'CA' and H'ca' are all equivalent.

Several legacy mpasm number formats are also supported. These formats have various shortcomings, but are still supported. The table below summarizes them.

base	general syntax	21 decimal written as
binary	[01]*b	10101b
octal	q'[0-7]*'	q'25'
octal	[0-7]*o	25o
octal	[0-7]*q	25q
decimal	0-9]*d	21d
decimal	.[0-9]*	.21
hex	[0-F]*h	15h

You can write the ASCII code for a character X using 'X', or A'X'.

3.2.4 Preprocessor

A line such as:

```
include foo.inc
```

will make gpasm fetch source lines from the file "foo.inc" until the end of the file, and then return to the original source file at the line following the include.

Lines beginning with a "#" are preprocessor directives, and are treated differently by gpasm. They may contain a "#define", or a "#undefine" directive.

Once gpasm has processed a line such as:

```
#define X Y
```

every subsequent occurrence of X is replaced with Y, until the end of file or a line

```
#undefine X
```

appears.

The preprocessor will replace an occurance of #v(expression) in a symbol with the value of "expression" in decimal. In the following expression:

```
number equ 5
label_#v( (number +1) * 5 )_suffix equ 0x10
```

gpasm will place the symbol "label_30_suffix" with a value of 0x10 in the symbol table.

The preprocessor in gpasm is only *like* the C preprocessor; its syntax is rather different from that of the C preprocessor. gpasm uses a simple internal preprocessor to implement "include", "#define" and "#undefine".

3.2.5 Processor header files

gputils distributes the Microchip processor header files. These files contain processor specific data that is helpful in developing PIC applications. The location of these files is reported in the gpasm help message. Use the INCLUDE directive to utilize the appropriate file in your source code. Only the name of the file is required. gpasm will search the default path automatically.

3.3 Directives

3.3.1 Code generation

In absolute mode, use the ORG directive to set the PIC memory location where gpasm will start assembling code. If you don't specify an address with ORG, gpasm assumes 0x0000. In relocatable mode, use the CODE directive.

3.3.2 Configuration

You can choose the fuse settings for your PIC implementation using the __CONFIG directive, so that the hex file set the fuses explicitly. Naturally you should make sure that these settings match your PIC hardware design.

The __MAXRAM and __BADRAM directives specify which RAM locations are legal. These directives are mostly used in processor-specific configuration files.

3.3.3 Conditional assembly

The IF, IFNDEF, IFDEF, ELSE and ENDIF directives enable you to assemble certain sections of code only if a condition is met. In themselves, they do not cause gpasm to emit any PIC code. The example in section 3.3.4 for demonstrates conditional assembly.

3.3.4 Macros

gpasm supports a simple macro scheme; you can define and use macros like this:

```
any     macro  parm
        movlw  parm
        endm
...
        any    33
```

A more useful example of some macros in use is:

```
; Shift  reg  left
slf     macro    reg
        clrc
        rlf      reg,f
endm

; Scale  W by "factor".    Result  in "reg",  W unchanged.
scale   macro    reg,  factor
        if (factor  == 1)
            movwf  reg                  ; 1 X is easy
        else
            scale    reg,  (factor  / 2) ; W * (factor  / 2)
            slf      reg,f               ; double  reg
```

```
            if ((factor & 1) == 1)     ; if lo-bit set ..
                addwf   reg,f          ; .. add W to reg
            endif
        endif
    endm
```

This recursive macro generates code to multiply W by a constant “factor”, and stores the result in “reg”. So writing:

```
        scale   tmp,D'10'
```

is the same as writing:

```
        movwf   tmp         ; tmp = W
        clrc
        rlf     tmp,f       ; tmp = 2 * W
        clrc
        rlf     tmp,f       ; tmp = 4 * W
        addwf   tmp,f       ; tmp = (4 * W) + W = 5 * W
        clrc
        rlf     tmp,f       ; tmp = 10 * W
```

3.3.5 $

$ expands to the address of the instruction currently being assembled. If it’s used in a context other than an instruction, such as a conditional, it expands to the address the next instruction would occupy, since the assembler’s idea of current address is incremented after an instruction is assembled. $ may be manipulated just like any other number:

```
$
$ + 1
$ - 2
```

and can be used as a shortcut for writing loops without labels.

```
LOOP:   BTFSS  flag,0x00
        GOTO   LOOP
        BTFSS  flag,0x00
        GOTO   $ - 1
```

3.3.6 Suggestions for structuring your code

Nested IF operations can quickly become confusing. Indentation is one way of making code clearer. Another way is to add braces on IF, ELSE and ENDIF, like this:

```
        IF (this)   ; {
          ...
        ELSE        ; }{
          ...
        ENDIF       ; }
```

After you've done this, you can use your text editor's show-matching-brace to check matching parts of the IF structure. In vi this command is "%", in emacs it's M-C-f and M-C-b.

3.3.7 Directive summary

__BADRAM

__BADRAM <expression> [, <expression]*

Instructs gpasm that it should generate an error if there is any use of the given RAM locations. Specify a range of addresses with <lo>-<hi>. See any processor-specific header file for an example.

See also: __MAXRAM

__CONFIG

__CONFIG <expression>

Sets the PIC processor's configuration fuses.

__IDLOCS

__IDLOCS <expression> or __IDLOCS <expression1>,<expression2>

Sets the PIC processor's identification locations. For 12 and 14 bit processors, the four id locations are set to the hexadecimal value of expression. For 18cxx devices idlocation expression1 is set to the hexadecimal value of expression2.

__MAXRAM

__MAXRAM <expression>

Instructs gpasm that an attempt to use any RAM location above the one specified should be treated as an error. See any processor specific header file for an example.

See also: __BADRAM

BANKISEL

BANKISEL <label>

This directive generates bank selecting code for indirect access of the address specified by <label>. The directive is not available for all devices. It is only available for 14 bit and 16 bit devices. For 14 bit devices, the bank selecting code will set/clear the IRP bit of the STATUS register. It will use MOVLB or MOVLR in 16 bit devices.

See also: BANKSEL, PAGESEL

BANKSEL

```
BANKSEL  <label>
```

This directive generates bank selecting code to set the bank to the bank containing <label>. The bank selecting code will set/clear bits in the FSR for 12 bit devices. It will set/clear bits in the STATUS register for 14 bit devices. It will use MOVLB or MOVLR in 16 bit devices. MOVLB will be used for enhanced 16 bit devices.

See also: BANKISEL, PAGESEL

CBLOCK

```
CBLOCK  [<expression>]
        <label>[:<increment>][,<label>[:incr          ement> ]]
ENDC
```

Marks the beginning of a block of constants <label>. gpasm allocates values for symbols in the block starting at the value <expression> given to CBLOCK. An optional <increment> value leaves space after the <label> before the next <label>.

See also: EQU

CODE

```
<label>  CODE  <expression>
```

Only for relocatable mode. Creates a new machine code section in the output object file. <label> specifies the name of the section. If <label> is not specified the default name ".code" will be used. <expression> is optional and specifies the absolute address of the section.

See also: IDATA, UDATA

CONSTANT

```
CONSTANT  <label>=<expression>      [, <label>=<expression>]*
```

Permanently assigns the value obtained by evaluating <expression> to thesymbol <label>. Similar to SET and VARIABLE, except it can not be changed once assigned.

See also: EQU, SET, VARIABLE

DA

```
<label>  DA <expression>    [, <expression]*
```

Stores Strings in program memory. The data is stored as one 14 bit word representing two 7 bit ASCII characters.

See also: DT

DATA

```
DATA <expression>  [, <expression]*
```

Generates the specified data.
See also: DA, DB, DE, DW

DB

```
<label>  DB <expression>  [, <expression]*
```

Declare data of one byte. The values are packed two per word.
See also: DA, DATA, DE, DW

DE

```
<label>  DE <expression>  [, <expression]*
```

Define EEPROM data. Each character in a string is stored in a separate word.
See also: DA, DATA, DB, DW

DT

```
DT <expression>  [, <expression]*
```

Generates the specified data as bytes in a sequence of RETLW instructions.
See also: DATA

DW

```
<label>  DW <expression>  [, <expression]*
```

Declare data of one word.
See also: DA, DATA, DB, DW

ELSE

```
ELSE
```

Marks the alternate section of a conditional assembly block.
See also: IF, IFDEF, IFNDEF, ELSE, ENDIF

END

```
END
```

Marks the end of the source file.

ENDC

ENDC

Marks the end of a CBLOCK.
See also: CBLOCK

ENDIF

ENDIF

Ends a conditional assembly block.
See also: IFDEF, IFNDEF, ELSE, ENDIF

ENDM

ENDM

Ends a macro definition.
See also: MACRO

ENDW

ENDW

Ends a while loop.
See also: WHILE

EQU

<label> EQU <expression>

Permanently assigns the value obtained by evaluating <expression> to the symbol <label>. Similar to SET and VARIABLE, except it can not be changed once assigned.
See also: CONSTANT, SET

ERROR

ERROR <string>

Issues an error message.
See also: MESSG

ERRORLEVEL

ERRORLEVEL {0 | 1 | 2 | +<msgnum> | -<msgnum>}[, ...]

Sets the types of messages that are printed.

Setting	Affect
0	Messages, warnings and errors printed.
1	Warnings and error printed.
2	Errors printed.
-<msgnum>	Inhibits the printing of message <msgnum>.
+<msgnum>	Enables the printing of message <msgnum>.

See also: LIST

EXTERN

EXTERN <symbol> [, <symbol>]*

Only for relocatable mode. Delcare a new symbol that is defined in another object file.
See also: GLOBAL

EXITM

EXITM

Immediately return from macro expansion during assembly.
See also: ENDM

EXPAND

EXPAND

Expand the macro in the listing file.
See also: ENDM

FILL

<label> FILL <expression>,<count>

Generates <count> occurrences of the program word or byte <expression>. If expression is enclosed by parentheses, expression is a line of assembly.
See also: DATA DW ORG

GLOBAL

GLOBAL <symbol> [, <symbol>]*

Only for relocatable mode. Delcare a symbol as global.
See also: GLOBAL

IDATA

<label> IDATA <expression>

Only for relocatable mode. Creates a new initialized data section in the output object file. <label> specifies the name of the section. If <label> is not specified the default name ".idata" will be used. <expression> is optional and specifies the absolute address of the section. Data memory is allocated and the initialization data is placed in ROM. The user must provide the code to load the data into memory.
See also: CODE, UDATA

IF

IF <expression>

Begin a conditional assembly block. If the value obtained by evaluating <expression> is true (i.e. non-zero), code up to the following ELSE or ENDIF is assembled. If the value is false (i.e. zero), code is not assembled until the corresponding ELSE or ENDIF.
See also: IFDEF, IFNDEF, ELSE, ENDIF

IFDEF

IFDEF <symbol>

Begin a conditional assembly block. If <symbol> appears in the symbol table, gpasm assembles the following code.
See also: IF, IFNDEF, ELSE, ENDIF

IFNDEF

IFNDEF <symbol>

Begin a conditional assembly block. If <symbol>does not appear in the symbol table, gpasm assembles the following code.
See also: IF, IFNDEF, ELSE, ENDIF

LIST

LIST <expression> [, <expression>] *

Enables output to the list (".lst") file. All arguments are interpreted as decimal regardless of the current radix setting. "list n=0" may be used to prevent page breaks in the code section of the list file. Other options are listed in the table below:

option	description
b=nnn	Sets the tab spaces
f=<format>	Set the hex file format. Can be inhx8m, inhx8s, inhx16, or inhx32.
mm=[ON|OFF]	Memory Map on or off
n=nnn	Sets the number of lines per page
p = <symbol>	Sets the current processor
pe = <symbol>	Sets the current processor and enables the 18xx extended mode
r= [oct | dec | hex]	Sets the radix
st = [ON | OFF]	Symbol table dump on or off
w=[0 | 1| 2]	Sets the message level.
x=[ON|OFF]	Macro expansion on or off

See also: NOLIST, RADIX, PROCESSOR

LOCAL

LOCAL <symbol>[[=<expression>], [<symbol>[=<expression>]]*]

Declares <symbol> as local to the macro that's currently being defined. This means that further occurrences of <symbol> in the macro definition refer to a local variable, with scope and lifetime limited to the execution of the macro.

See also: MACRO, ENDM

MACRO

<label> MACRO [<symbol> [, <symbol>]*]

Declares a macro with name <label>. gpasm replaces any occurrences of <symbol> in the macro definition with the parameters given at macro invocation.

See also: LOCAL, ENDM

MESSG

MESSG <string>

Writes <string> to the list file, and to the standard error output.

See also: ERROR

NOEXPAND

NOEXPAND

Turn off macro expansion in the list file.

See also: EXPAND

NOLIST

NOLIST

Disables list file output.
See also: LIST

ORG

ORG <expression>

Sets the location at which instructions will be placed. If the source file does not specify an address with ORG, gpasm assumes an ORG of zero.

PAGE

PAGE

Causes the list file to advance to the next page.
See also: LIST

PAGESEL

PAGESEL <label>

This directive will generate page selecting code to set the page bits to the page containing the designated <label>. The page selecting code will set/clear bits in the STATUS for 12 bit and 14 bit devices. For 16 bit devices, it will generate MOVLW and MOVWF to modify PCLATH. The directive is ignored for enhanced 16 bit devices.
See also: BANKISEL, BANKSEL

PROCESSOR

PROCESSOR <symbol>

Selects the target processor. See section **??** for more details.
See also: LIST

RADIX

RADIX <symbol>

Selects the default radix from “oct” for octal, “dec” for decimal or “hex” for hexadecimal. gpasm uses this radix to interpret numbers that don’t have an explicit radix.
See also: LIST

RES

RES <mem_units>

Causes the memory location pointer to be advanced <mem_units>. Can be used to reserve data storage.
See also: FILL, ORG

SET

<label> SET <expression>

Temporarily assigns the value obtained by evaluating <expression> to the symbol <label>.
See also: SET

SPACE

SPACE <expression>

Inserts <expression> number of blank lines into the listing file.
See also: LIST

SUBTITLE

SUBTITLE <string>

This directive establishes a second program header line for use as a subtitle in the listing output. <string> is an ASCII string enclosed by double quotes, no longer than 60 characters.
See also: TITLE

TITLE

TITLE <string>

This directive establishes a program header line for use as a title in the listing output. <string> is an ASCII string enclosed by double quotes, no longer than 60 characters.
See also: SUBTITLE

UDATA

<label> UDATA <expression>

Only for relocatable mode. Creates a new uninitialized data section in the output object file. <label> specifies the name of the section. If <label> is not specified the default name ".udata" will be used. <expression> is optional and specifies the absolute address of the section.
See also: CODE, IDATA, UDATA_ACS, UDATA_OVR, UDATA_SHR

UDATA_ACS

<label> UDATA_ACS <expression>

Only for relocatable mode. Creates a new uninitialized accessbank data section in the output object file. <label> specifies the name of the section. If <label> is not specified the default name ".udata_acs" will be used. <expression> is optional and specifies the absolute address of the section.

See also: CODE, IDATA, UDATA

UDATA_OVR

<label> UDATA_OVR <expression>

Only for relocatable mode. Creates a new uninitialized overlaid data section in the output object file. <label> specifies the name of the section. If <label> is not specified the default name ".udata_ovr" will be used. <expression> is optional and specifies the absolute address of the section.

See also: CODE, IDATA, UDATA

UDATA_SHR

<label> UDATA_SHR <expression>

Only for relocatable mode. Creates a new uninitialized sharebank data section in the output object file. <label> specifies the name of the section. If <label> is not specified the default name ".udata_shr" will be used. <expression> is optional and specifies the absolute address of the section.

See also: CODE, IDATA, UDATA

VARIABLE

VARIABLE <label>[=<expression>, <label>[=<expression>]]*

Delcares variable with the name <label>. The value of <label> may later be reassigned. The value of <label> does not have to be assigned at declaration.

See also: CONSTANT

WHILE

WHILE <expression>

Performs loop while <expression> is true.

See also: ENDW

3.3.8 High level extensions

gpasm supports several directives for use with high level languages. These directives are easily identified because they start with ".". They are only available in relocatable mode.

These features are advanced and require knowledge of how gputils relocatable objects work. These features are intended to be used by compilers. Nothing prevents them from being used with assembly.

.DEF

.DEF <symbol> [, <expression>]*

Create a new COFF <symbol>. Options are listed in the table below:

option	description
absolute	Absolute symbol keyword
class=nnn	Sets the symbol class (byte sized)
debug	Debug symbol keyword
extern	External symbol keyword
global	Global symbol keyword
size=nnn	Reserve words or bytes for the symbol
static	Static Symbol keyword
type=nnn	Sets the symbol type (short sized)
value=nnn	Sets the symbol value

This directive gives the user good control of the symbol table. This control is necessary, but if used incorrectly it can have many undesirable consequences. It can easily cause errors during linking or invalid machine code. The user must fully understand the operation of gputils COFF symbol table before modifying its contents.

For best results, only one of the single keywords should be used. The keyword should follow the symbol name. The keyword should then be followed by any expressions that directly set the values. Here is an example:

.def global_clock, global, type = T_ULONG, size = 4

See also: .DIM

.DIM

.DIM <symbol>, <number>, <expression> [, <expression>] *

Create <number> auxiliary symbols, attached to <symbol>. Fill the auxilary symbols with the values specified in <expression>. The expressions must result in byte sized values when evaluated or be strings. The symbol must be a COFF symbol.

This directive will generate an error if the symbol already has auxiliary symbols. This prevents the user from corrupting automatically generated symbols.

Each auxiliary symbol is 18 bytes. So the contents specified by the expressions must be less than or equal to 18 * <number>.

gpasm does not use auxilary symbols. So the contents have no effect on its operation. However, the contents may be used by gplink or a third party tool.

See also: .DEF

.DIRECT

.DIRECT <command>, <string>

Provides a mechanism for direct communication from the program to the debugging environment. This method has no impact on the executable. The symbols will appear in both the COFF files and the COD files.

Each directive creates a new COFF symbol ".direct". An auxilary symbol is attached that contains <command> and <string>. The string must be less than 256 bytes. The command must have a value 0 to 255. There are no restrictions on the content, however these messages must conform to the debugging environment. The typical values are summarized in the table below:

ASCII command	description
a	User defined assert
A	Assembler/Compiler defined assert
e	User defined emulator commands
E	Assembler/Compiler defined emulator commands
f	User defined printf
F	Assembler/Compiler defined printf
l	User defined log command
L	Assembler/Compiler/Code verification generated log command

The symbols also contain the address where the message was inserted into the assembly. The symbols, with the final relocated addresses, are available in executable COFF. The symbols are also written to the COD file. They can be viewed using gpvc.

See also: .DEF, .DIM

.EOF

.EOF

This directive causes an end of file symbol to be placed in the symbol table. Normally this symbol is automatically generated. This directive allows the user to manually generate the symbol. The directive is only processed if the "-g" command line option is used. When that option is used, the automatic symbol generation is disabled.

See also: .EOF, .FILE, .LINE

.FILE

.FILE <string>

This directive causes a file symbol to be placed in the symbol table. Normally this symbol is automatically generated. This directive allows the user to manually generate the symbol. The directive is only processed if the "-g" command line option is used. When that option is used, the automatic symbol generation is disabled.

See also: .EOF, .FILE, .LINE

.IDENT

.IDENT <string>

Creates an .ident COFF symbol and appends an auxiliary symbol. The auxiliary symbol points to an entry in the string table. The entry contains <string>. It is an ASCII comment of any length. This symbol has no impact on the operation of gputils. It is commonly used to store compiler versions.

See also: .DEF, .DIM

.LINE

.LINE <expression>

This directive causes and COFF line number to be generated. Normally they are automatically generated. This directive allows the user to manually generate the line numbers. The directive is only processed if the "-g" command line option is used. When that option is used, the automatic symbol generation is disabled. The <expression> is always evaluated as decimal regardless of the current radix setting.

See also: .EOF, .FILE, .LINE

.TYPE

.TYPE <symbol>, <expression>

This directive modifies the COFF type of an existing <symbol>. The symbol must be defined. The type must be 0 to 0xffff. Common types are defined in coff.inc.

COFF symbol types default to NULL in gpasm. Although the type has no impact linking or generating an executable, it does help in the debug environment.

See also: .DEF

3.4 Instructions

3.4.1 Instruction set summary

12 bit Devices (PIC12C5XX)

Syntax	Description
ADDLW <imm8>	Add immediate to W
ADDWF <f>,<dst>	Add W to <f>, result in <dst>
ANDLW <imm8>	And W and literal, result in W
ANDWF <f>,<dst>	And W and <f>, result in <dst>
BCF <f>,<bit>	Clear <bit> of <f>
BSF <f>,<bit>	Set <bit> of <f>
BTFSC <f>,<bit>	Skip next instruction if <bit> of <f> is clear
BTFSS <f>,<bit>	Skip next instruction if <bit> of <f> is set
CALL <addr>	Call subroutine
CLRF <f>,<dst>	Write zero to <dst>
CLRW	Write zero to W
CLRWDT	Reset watchdog timer
COMF <f>,<dst>	Complement <f>, result in <dst>
DECF <f>,<dst>	Decrement <f>, result in <dst>
DECFSZ <f>,<dst>	Decrement <f>, result in <dst>, skip if zero
GOTO <addr>	Go to <addr>
INCF <f>,<dst>	Increment <f>, result in <dst>
INCFSZ <f>,<dst>	Increment <f>, result in <dst>, skip if zero
IORLW <imm8>	Or W and immediate
IORWF <f>,<dst>	Or W and <f>, result in <dst>
MOVF <f>,<dst>	Move <f> to <dst>
MOVLW <imm8>	Move literal to W
MOVWF <f>	Move W to <f>
NOP	No operation
OPTION	
RETLW <imm8>	Load W with immediate and return
RLF <f>,<dst>	Rotate <f> left, result in <dst>
RRF <f>,<dst>	Rotate <f> right, result in <dst>
SLEEP	Enter sleep mode
SUBWF <f>,<dst>	Subtract W from <f>, result in <dst>
SWAPF <f>,<dst>	Swap nibbles of <f>, result in <dst>
TRIS	
XORLW	Xor W and immediate
XORWF	Xor W and <f>, result in <dst>

14 Bit Devices (PIC16CXX)

Syntax	Description
ADDLW <imm8>	Add immediate to W
ADDWF <f>,<dst>	Add W to <f>, result in <dst>
ANDLW <imm8>	And immediate to W
ANDWF <f>,<dst>	And W and <f>, result in <dst>
BCF <f>,<bit>	Clear <bit> of <f>
BSF <f>,<bit>	Set <bit> of <f>
BTFSC <f>,<bit>	Skip next instruction if <bit> of <f> is clear
BTFSS <f>,<bit>	Skip next instruction if <bit> of <f> is set
CALL <addr>	Call subroutine
CLRF <f>,<dst>	Write zero to <dst>
CLRW	Write zero to W
CLRWDT	Reset watchdog timer
COMF <f>,<dst>	Complement <f>, result in <dst>
DECF <f>,<dst>	Decrement <f>, result in <dst>
DECFSZ <f>,<dst>	Decrement <f>, result in <dst>, skip if zero
GOTO <addr>	Go to <addr>
INCF <f>,<dst>	Increment <f>, result in <dst>
INCFSZ <f>,<dst>	Increment <f>, result in <dst>, skip if zero
IORLW <imm8>	Or W and immediate
IORWF <f>,<dst>	Or W and <f>, result in <dst>
MOVF <f>,<dst>	Move <f> to <dst>
MOVLW <imm8>	Move literal to W
MOVWF <f>	Move W to <f>
NOP	No operation
OPTION	
RETFIE	Return from interrupt
RETLW <imm8>	Load W with immediate and return
RETURN	Return from subroutine
RLF <f>,<dst>	Rotate <f> left, result in <dst>
RRF <f>,<dst>	Rotate <f> right, result in <dst>
SLEEP	Enter sleep mode
SUBLW	Subtract W from literal
SUBWF <f>,<dst>	Subtract W from <f>, result in <dst>
SWAPF <f>,<dst>	Swap nibbles of <f>, result in <dst>
TRIS	
XORLW	Xor W and immediate
XORWF	Xor W and <f>, result in <dst>

Ubicom Processors

For Ubicom (Scenix) processors, the assembler supports the following instructions, in addition to those listed under "12 Bit Devices" above.

Syntax	Description
BANK <imm3>	
IREAD	
MODE <imm4>	
MOVMW	
MOVWM	
PAGE <imm3>	
RETI	
RETIW	
RETP	
RETURN	

Special macros

There are also a number of standard additional macros. These macros are:

Syntax	Description
ADDCF <f>,<dst>	Add carry to <f>, result in <dst>
B <addr>	Branch
BC <addr>	Branch on carry
BZ <addr>	Branch on zero
BNC <addr>	Branch on no carry
BNZ <addr>	Branch on not zero
CLRC	Clear carry
CLRZ	Clear zero
SETC	Set carry
SETZ	Set zero
MOVFW <f>	Move file to W
NEGF <f>	Negate <f>
SKPC	Skip on carry
SKPZ	Skip on zero
SKPNC	Skip on no carry
SKPNZ	Skip on not zero
SUBCF <f>,<dst>	Subtract carry from <f>, result in <dst>
TSTF <f>	Test <f>

3.5 Errors/Warnings/Messages

gpasm writes every error message to two locations:

- the standard error output
- the list file (".lst")

The format of error messages is:

Error <src-file> <line> : <code> <description>

where:

<src-file> is the source file where gpasm encountered the error

<line> is the line number

<code> is the 3-digit code for the error, given in the list below

<description> is a short description of the error. In some cases this contains further information about the error.

Error messages are suitable for parsing by emacs' "compilation mode". This chapter lists the error messages that gpasm produces.

3.5.1 Errors

101 ERROR directive

A user-generated error. See the ERROR directive for more details.

114 Divide by zero

gpasm encountered a divide by zero.

115 Duplicate Label

Duplicate label or redefining a symbol that can not be redefined.

124 Illegal Argument

gpasm encountered an illegal argument in an expression.

125 Illegal Condition

An illegal condition like a missing ENDIF or ENDW has been encountered.

126 Argument out of Range

The expression has an argument that was out of range.

127 Too many arguments

gpasm encountered an expression with too many arguments.

128 Missing argument(s)

gpasm encountered an expression with at least one missing argument.

129 Expected

Expected a certain type of argument.

130 Processor type previously defined

The processor is being redefined.

131 Undefined processor

The processor type has not been defined.

132 Unknown processor

The selected processor is not valid. Check the processors listed in section **??**.

133 Hex file format INHX32 required

An address above 32K was specified.

135 Macro name missing

A macro was defined without a name.

136 Duplicate macro name

A macro name was duplicated.

145 Unmatched ENDM

ENDM found without a macro definition.

159 Odd number of FILL bytes

In PIC18CXX devices the number of bytes must be even.

3.5.2 Warnings

201 Symbol not previously defined.

The symbol being #undefined was not previously defined.

202 Argument out of range

The argument does not fit in the allocated space.

211 Extraneous arguments

Extra arguments were found on the line.

215 Processor superseded by command line

The processor was specified on the command line and in the source file. The command line has precedence.

216 Radix superseded by command line

The radix was specified on the command line and in the source file. The command line has precedence.

217 Hex format superseded by command line

The hex file format was specified on the command line and in the source file. The command line has precedence.

218 Expected DEC, OCT, HEX. Will use HEX.

gpasm encountered an invalid radix.

219 Invalid RAM location specified.

gpasm encountered an invalid RAM location as specified by the __MAXRAM and __BADRAM directives.

222 Error messages can not be disabled

Error messages can not be disabled using the ERRORLEVEL directive.

223 Redefining processor

The processor is being reselected by the LIST or PROCESSOR directive.

224 Use of this instruction is not recommanded

Use of the TRIS and OPTION instructions is not recommended for a PIC16CXX device.

3.5.3 Messages

301 User Message

User message, invoked with the MESSG directive.

303 Program word too large. Truncated to core size.

gpasm has encounter a program word larger than the core size of the selected device.

304 ID Locations value too large. Last four hex digits used.

The ID locations value specified is too large.

305 Using default destination of 1 (file).

No destination was specified so the default location was used.

308 Warning level superseded by command line

The warning level was specified on the command line and in the source file. The command line has precedence.

309 Macro expansion superseded by command line

Macro expansion was specified on the command line and in the source file. The command line has precedence.

Chapter 4

gplink

gplink relocates and links gpasm COFF objects and generates an absolute executable COFF.

4.1 Running gplink

The general syntax for running gplink is

```
gplink [options] [objects] [libraries]
```

Where options can be one of:

Option	Meaning
a	Produce hex file in one of four formats: inhx8m, inhx8s, inhx16, inhx32 (the default)
c	Output an executable object
d	Display debug messages
f <value>	Fill unused unprotected program memory with <value>
h	Show the help message
I <directory>	Specify an include directory
l	Disable the list file output
m	Output a map file
o <file>	Alternate name of hex output file
q	Quiet
r	Attempt to relocate unshared data sections to shared memory if relocation fails
s <file>	Specify linker script
-t <value>	Create a stack section
v	Print gplink version information and exit

4.2 gplink outputs

gplink creates an absolute executable COFF. From this COFF a hex file and cod file are created. The executable COFF is only written when the "-c" option is added. This file is useful for simulating the design with mpsim. The cod file is used for simulating with gpsim.

gplink can also create a map file. The map file reports the final addresses gplink has assigned to the COFF sections. This is the same data that can be viewed in the executable COFF with gpvo.

4.3 Linker scripts

gplink requires a linker script. This script tells gplink what memory is available in the target processor. A set of Microchip generated scripts are installed with gputils. These scripts were intended as a starting point, but for many applications they will work as is.

If the user does not specify a linker script, gplink will attempt to use the default script for the processor reported in the object file. The default location of the scripts is reported in the gplink help message.

4.4 Stacks

gplink can create a stack section at link time using a stack directive in the linker script. The same feature can be utilized with a -t option on the command line. gplink will create the section and two symbols. _stack points to the beginning of the stack section and _stack_end points to the end.

Chapter 5

gplib

gplib creates, modifies and extracts COFF archives. This allows a related group of objects to be combined into one file. Then this one file is passed to gplink.

5.1 Running gplib

The general syntax for running gplib is

```
gplib [options] library [member]
```

Where options can be one of:

Option	Meaning
c	Create a new library
d	Delete member from library
h	Show the help message
n	Don't add the symbol index
q	Quiet mode
r	Add or replace member from library
s	List global symbols in libary
t	List member in library
v	Print gplib version information and exit
x	Extract member from library

5.2 Creating an archive

The most common operation is to create a new archive:

```
gplib -c math.a mult.o add.o sub.o
```

This command will create a new archive “math.a” that contains “mult.o add.o sub.o”.

The name of the archive “math.a” is arbitrary. The tools do not use the file extension to determine file type. It could just as easily been “math.lib” or “math”.

When you use the library, simply add it to the list of object passed to gplink. gplink will scan the library and only extract the archive members that are required to resolve external references. So the application won’t necessarily contain the code of all the archive members.

5.3 Other gplib operations

Most of the other are useful , but will be used much less often. For example you can replace individual archive members, but most people elect to delete the old archive and create a new one.

5.4 Archive format

The file format is a standard COFF archive. A header is added to each member and the unmodified object is copied into the archive.

Being a standard archive they do include a symbol index. It provides a simple why to determine which member should be extract to resolve external references. This index is not included in mplib archives. So using gplib archives with Microchip Tools will probably cause problems unles the “-n” option is added when the archive is created.

Chapter 6

Utilities

6.1 gpdasm

gpdasm is a disassembler for gputils. It converts hex files generated by gpasm and gplink into disassembled instructions.

6.1.1 Running gpdasm

The general syntax for running gpdasm is

```
gpdasm [options] hex-file
```

Where options can be one of:

Option	Meaning
c	Decode special mnemonics
h	Display the help message
i	Display hex file information
l	List supported processors
m	Memory dump hex file
p<processor>	Select processor
s	Print short form output
v	Print gpdasm version information and exit
y	Enable 18xx extended mode

gpdasm doesn't specifically create an output file. It dumps its output to the screen. This helps to reduce the risk that a good source file will be unintentionally overwritten. If you want to create an output file and your shell is "sh", "bash" or "ksh", you can do something like:

```
gpdasm test.hex > test.dis
```

This redirects standard output to the file "test.dis".

6.1.2 Comments on Disassembling

- The gpdasm only uses a hex file as an input. Because of this it has no way to distinguish between instructions and data in program memory.
- If gpdasm encounters an unknown instruction it uses the DW directive and treats it as raw data.
- There are DON'T CARE bits in the instruction words. Normally, this isn't a problem. It could be, however, if a file with data in the program memory space is disassembled and then reassembled. Example: gpdasm will treat 0x0060 in a 14 bit device as a NOP. If the output is then reassembled, gpasm will assign a 0x0000 value. The value has changed and both tools are behaving correctly.

6.2 gpvc

gpvc is cod file viewer for gputils. It provides an easy way to view the contents of the cod files generated by gpasm and gplink.

6.2.1 Running gpvc

The general syntax for running gpvc is

```
gpvc [options] cod-file
```

Where options can be one of:

Option	Meaning
a	Display all information
d	Display directory header
s	Display symbols
h	Show the help message.
r	Display ROM
l	Display source listing
m	Display debug message area
v	Print gpvc version information and exit

gpvc doesn't specifically create an output file. It dumps its output to the screen. If you want to create an output file and your shell is "sh", "bash" or "ksh", you can do something like:

```
gpvc test.cod > test.dump
```

This redirects standard output to the file "test.dump".

6.3 gpvo

gpvo is COFF object file viewer for gputils. It provides an easy way to view the contents of objects generated by gpasm and gplink.

6.3.1 Running gpvo

The general syntax for running gpvo is

```
gpvo [options] object-file
```

Where options can be one of:

Option	Meaning
b	Binary data
c	Decode special mnemonics
f	File header
h	Show the help message
n	Suppress filenames
s	Section data
t	Symbol data
v	Print gpvo version information and exit
y	Enable 18xx extended mode

gpvo doesn't specifically create an output file. It dumps its output to the screen. If you want to create an output file and your shell is "sh", "bash" or "ksh", you can do something like:

```
gpvo test.obj > test.dump
```

This redirects standard output to the file "test.dump".

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